

Lecture 4: Chernoff/Concentration, PAC-learning (Sept. 14)

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4.1 Outline

1. Concentration inequalities:
Chernoff's inequality, multiplicative Chernoff's inequality; Bennett's inequality, Bernstein inequality
2. PAC-learning:
PAC learnability based on 'fitness'/union bounds

4.2 Concentration inequalities

Theorem 4.1 (Chernoff's Inequality). Let $X_1, \dots, X_n \in [0, 1]$ be i.i.d. random variables, $\bar{X}_n = \frac{1}{n}(X_1 + \dots + X_n)$, $\mu = \mathbb{E}X_1$. We have

(a) $\forall \delta \in (0, 1)$, with probability $1 - \delta$,

$$\bar{X}_n \leq \mu + \sqrt{\frac{\log(1/\delta)}{2n}};$$

(b) $\forall \delta \in (0, 1)$, with probability $1 - \delta$,

$$\bar{X}_n \geq \mu - \sqrt{\frac{\log(1/\delta)}{2n}}.$$

Proof. Since $X_1 \in [a, b]$ implies that X_1 is $\sigma(X_1)$ -SG with $\sigma(X_1) = \frac{b-a}{n}$, $X_1 \in [0, 1]$ indicates that

$$\sigma(\bar{X}_n) = \frac{\sigma(X_1)}{\sqrt{n}} = \frac{1}{2\sqrt{n}}.$$

Applying this fact with Hoeffding inequality, the Chernoff's inequality is proven. \square

Theorem 4.2 (Multiplicative Chernoff's Inequality). Let $X_1, \dots, X_n \in [0, 1]$ be i.i.d. random variables, $\bar{X}_n = \frac{1}{n}(X_1 + \dots + X_n)$, $\mu = \mathbb{E}X_1$. We have

(a) $\forall \delta \in (0, 1)$, with probability $1 - \delta$,

$$\bar{X}_n \leq \mu + \sqrt{\frac{2\mu \log(1/\delta)}{n}} + \frac{1}{3n};$$

(b) $\forall \delta \in (0, 1)$, with probability $1 - \delta$,

$$\bar{X}_n \geq \mu - \sqrt{\frac{2\mu \log(1/\delta)}{n}}. \quad (*)$$

Remark 4.3.

(a) How big can μ be?

By (*): $\mu \leq \bar{X}_n + \sqrt{\frac{2\mu \log(1/\delta)}{n}}$.

(b) Let

$$f(a, c) = \max\{u : u \leq a + \sqrt{u \cdot c}\}, \text{ where } a = \bar{X}_n, c = \frac{2 \log(1/\delta)}{n}.$$

Then

$$\begin{aligned} \mu + \frac{\log(1/\delta)}{2n} &\geq \sqrt{\frac{2\mu \log(1/\delta)}{n}} \text{ and equality holds when } \mu = \frac{\log(1/\delta)}{2n}, \\ \Rightarrow \inf_{0 < \gamma < 1} \gamma\mu + \frac{\log(1/\delta)}{2\gamma n} &= \sqrt{\frac{2\mu \log(1/\delta)}{n}}, \\ \Rightarrow \mu - \sqrt{\frac{2\mu \log(1/\delta)}{n}} &= \sup_{0 < \gamma < 1} (1 - \gamma)\mu - \frac{\log(1/\delta)}{2\gamma n}. \end{aligned}$$

Let $\gamma = 1/2$, then with (*) we have

$$\bar{X}_n \geq \frac{\mu}{2} - \frac{\log(1/\delta)}{n}.$$

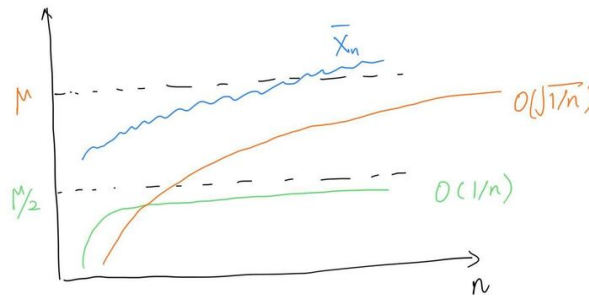


Figure 4.1: Example: set $\gamma = 1/2$.

(c) When we apply γ that does not maximize the term $(1 - \gamma)\mu - \frac{\log(1/\delta)}{2\gamma n}$, we cannot claim that we get a better ‘convergence’ rate, because when $n \rightarrow \infty$, $(1 - \gamma)\mu - \frac{\log(1/\delta)}{2\gamma n}$ and \bar{X}_n converges to different values. In detail, \bar{X}_n converges to μ regardless of the value of γ , and $(1 - \gamma)\mu - \frac{\log(1/\delta)}{2\gamma n}$ converges to $(1 - \gamma)\mu \neq \mu$ when $0 < \gamma < 1$.

To say something about the convergence of \bar{X}_n , we need to have the coefficient of μ be 1.

Theorem 4.4 (Bernett’s Inequality). *Let X_1, \dots, X_n be i.i.d. random variables. Set $\bar{X}_n = \frac{1}{n}(X_1 + \dots + X_n)$ and $\mu = \mathbb{E}X_1$. If $X_1 - \mu \leq b$, with probability $1 - \delta$, we have*

$$\bar{X}_n \leq \mu + \sqrt{\frac{2 \text{Var}(X_1) \log(1/\delta)}{n}} + \frac{b}{3n}.$$

4.3 PAC-learning (L. Valiant)

Let function $f_* : \{0, 1\}^d \rightarrow \{0, 1\}$, $X_1, X_2, \dots, X_n \in \{0, 1\}^d := \underline{2}^d$ be i.i.d. random variables drawn from distribution P_X , data set $D_n = \{(X_1, f_*(X_1)), \dots, (X_n, f_*(X_n))\}$.

Let $f_* \in \mathcal{F} \subset \underline{2}^{2^d}$ and $f \in \underline{2}^{2^d}$, $P_X^{f_*} := P(X_1, f_* X_1)$, and

$$\begin{aligned} L(f) &= \mathbb{P}(f(X) \neq f_*(X)) = L(P_X^{f_*}, f), \\ l : \underline{2} \times \underline{2} &\rightarrow \underline{2}, \quad l(y, y') = \mathbf{1}(y \neq y'), \\ L(P_X^{f_*}, f) &= \int P(dx, dy) l(f(x), y). \end{aligned}$$

Definition 4.5 (PAC-Learning). Fix $\mathcal{C} = (\mathcal{C}_d)_{d \geq 1}$, where $\mathcal{C}_d \subset \underline{2}^{2^d}$. \mathcal{C} is **PAC-learnable (Probably Approximately Correctly)** if \exists polynomial $p \in \mathbb{R}[x, y, z]$ and $\mathcal{A} = (\mathcal{A}_{n,d})_{n \geq 1, d \geq 1}$ where $\mathcal{A}_{n,d} : (\underline{2}^d \times \underline{2})^n \rightarrow \underline{2}^{2^d}$

$$\begin{aligned} \text{s.t. } \forall \varepsilon \in (0, 1), \delta \in (0, 1), d \geq 1, P \in \mathcal{M}_1(\underline{2}^d), f_* \in \mathcal{C}_d, \\ n \geq \lceil p(1/\varepsilon, 1/\delta, d) \rceil, \\ X_1, X_2, \dots, X_n \sim P_X, \\ f_n = \mathcal{A}_{n,d} \left(\underbrace{(X_1, f_*(X_1)), \dots, (X_n, f_*(X_n))}_{D_n} \right), \end{aligned}$$

we have

$$\mathbb{P} \left(L \left(P_X^{f_*}, f_n \right) \geq \varepsilon \right) \leq \delta.$$

In other words, with probability $1 - \delta$, $\mathbb{P}(f_n(X) \neq f_*(x) | D_n) \leq \varepsilon$.

Remark 4.6. (a) Example:

$$\mathcal{C}_{\text{AND}}^d = \left\{ f : \underline{2}^d \rightarrow \underline{2} \mid \exists u \subset [d], \forall x \in \underline{2}^d : f(x) = \min_{j \in u} X_j \right\}.$$

(b) (i) $L(f_*) = 0$. (ii) When $Y_i = f_*(X_i)$, there is **NO** noise and this will make learning **faster**.